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# Solutions Manual

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# Chapter 10

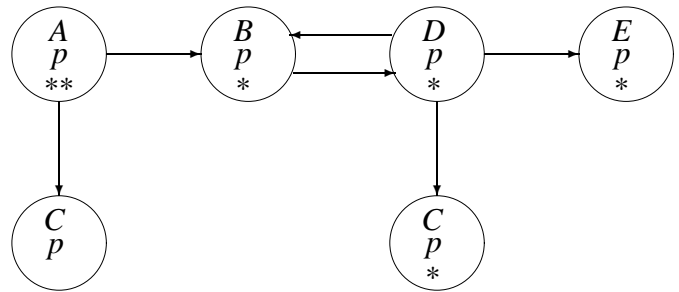
## Section 10.1

### Exercise 10.1.1

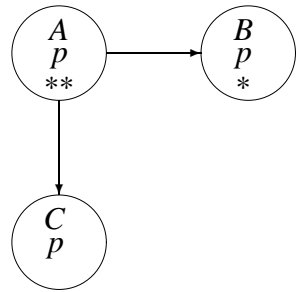
- (a) SELECT on MovieStar, SELECT on MovieExec.
- (b) SELECT on MovieExec, SELECT on Movies, SELECT on StarsIn.
- (c) SELECT on Movies, SELECT on Studio, INSERT on Studio (or INSERT(name) on Studio).
- (d) DELETE on StarsIn.
- (e) UPDATE on MovieExec (or UPDATE(name) on MovieExec).
- (f) REFERENCES on MovieStar (or REFERENCES(gender, name) on MovieStar).
- (g) REFERENCES on Studio, REFERENCES on MovieExec (or REFERENCES(name, presC#) on Studio, REFERENCES(cert#, netWorth) on MovieExec).

**Exercise 10.1.2**

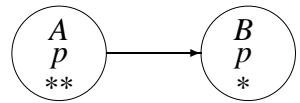
After step (4), the grant diagram is as follows:



After step (5), the grant diagram is as follows:

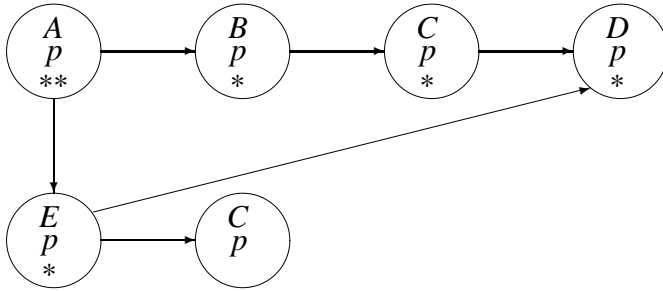


After step (6), the grant diagram is as follows:

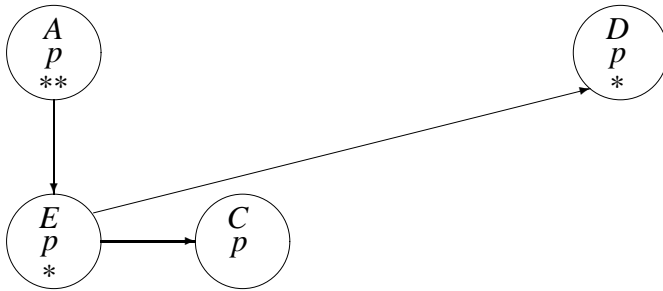


### Exercise 10.1.3

After step (5), the grant diagram is as follows:



After step (6), the grant diagram is as follows:



### Exercise 10.1.4

The grant diagram after the final step is as follows:



## Section 10.2

### Exercise 10.2.1

(a) The rules for trips that have reasonable connections are:

$$\begin{aligned}\text{Trips}(x, y, \text{dep}, \text{arr}) &\leftarrow \text{Flights}(\_, x, y, \text{dep}, \text{arr}) \\ \text{Trips}(x, y, \text{dep}, \text{arr}) &\leftarrow \text{Trips}(x, z, \text{dep1}, \text{arr1}) \text{ AND} \\ &\quad \text{Trips}(z, y, \text{dep2}, \text{arr2}) \text{ AND} \\ &\quad \text{arr1} \leq \text{dep2} - 100\end{aligned}$$

(b) Using the book's syntax, the SQL is:

```
WITH RECURSIVE Trips(frm, to, dep, arr) AS
  (SELECT frm, to, dep, arr
   FROM   Flights
   )
UNION
  (SELECT T.frm, F.to, T.dep, F.arr
   FROM   Trips T, Flights F
   WHERE  T.to = F.from
        AND T.arr <= F.dep - 100
   )
SELECT *
FROM   Trips;
```

### Exercise 10.2.2

Because `FROM` is one of the SQL reserved words, using it as an identifier is not recommended. Note that most major vendors do not prohibit the use of reserved words when the use is not ambiguous (e.g. `SELECT FROM FROM FROM` is not ambiguous and will work), but such use is highly discouraged for readability and portability reasons.

### Exercise 10.2.3

(a)

$$\begin{aligned}\text{FollowOn}(x, y) &\leftarrow \text{SequelOf}(x, y) \\ \text{FollowOn}(x, y) &\leftarrow \text{FollowOn}(x, z) \text{ AND} \\ &\quad \text{SequelOf}(z, y)\end{aligned}$$

(b) Using the book's syntax, the SQL is:

```
WITH RECURSIVE FollowOn(movie, followOn) AS
  (SELECT movie, sequel
   FROM SequelOf )
UNION
  (SELECT F.movie, S.sequel
   FROM FollowOn F, Sequel S
   WHERE F.followOn = S.movie)
SELECT *
FROM FollowOn;
```

(c) Using the book's syntax, the SQL is:

```
WITH RECURSIVE FollowOn(movie, followOn) AS
  (SELECT movie, sequel
   FROM SequelOf )
UNION
  (SELECT F.movie, S.sequel
   FROM FollowOn F, Sequel S
   WHERE F.followOn = S.movie)
SELECT movie, followOn
FROM FollowOn
EXCEPT
SELECT movie, sequel
FROM SequelOf;
```

(Similarly, NOT IN or NOT EXISTS can be used instead of EXCEPT).

- (d) One of the ways is to first get all of the recursive tuples as for the original FollowOn in (a), and then subtract the those tuples that represent sequel or sequel of a sequel. Using the book's syntax, the SQL would be:

```
WITH RECURSIVE FollowOn(movie, followOn) AS
  (SELECT movie, sequel
   FROM   SequelOf      )
  UNION
  (SELECT F.movie, S.sequel
   FROM   FollowOn F, Sequel S
   WHERE  F.followOn = S.movie)
SELECT movie, followOn
FROM   FollowOn
EXCEPT
(SELECT movie, sequel
 FROM   SequelOf
 UNION
 SELECT X.movie, Y.sequel
 FROM   Sequel X, Sequel Y
 WHERE  X.sequel = Y.movie);
```

Another way would be to start FollowOn tuples only from the tuples of movies that have more than two sequels (using a join similar to the one above but with three Sequel tables).

- (e) We simply need to count the number of followon values per movie. Using the book's syntax, the SQL would be:

```
WITH RECURSIVE FollowOn(movie, followOn) AS
  (SELECT movie, sequel
   FROM   SequelOf      )
  UNION
  (SELECT F.movie, S.sequel
   FROM   FollowOn F, Sequel S
   WHERE  F.followOn = S.movie)
SELECT movie
```

```

FROM FollowOn
GROUP BY movie
HAVING COUNT(followon) >= 2;

```

- (f) This is, in a sense, a reverse of (e) above, because to have at most one followon means that the total count of the tuples grouped by the given movie  $x$  must be no greater than 2 (one for the movie and its sequel, and the other for the sequel and its sequel). Using the book's syntax, the SQL would be:

```

WITH RECURSIVE FollowOn(movie, followon) AS
  (SELECT movie, sequel
   FROM SequelOf )
  UNION
  (SELECT F.movie, S.sequel
   FROM FollowOn F, Sequel S
   WHERE F.followon = S.movie)
SELECT movie, followon
FROM FollowOn
WHERE movie IN(SELECT movie
                FROM FollowOn
                GROUP BY movie
                HAVING COUNT(followon) <= 2);

```

### Exercise 10.2.4

- (a) WITH RECURSIVE Path(class, rclass) AS  
 (SELECT class, rclass  
 FROM Rel )  
 UNION  
 (SELECT Path.class, Rel.rclass  
 FROM Path, Rel  
 WHERE Path.rclass = Rel.class)  
 SELECT \*  
 FROM Path;
- (b) WITH RECURSIVE Path(class, rclass) AS  
 (SELECT class, rclass



```

        FROM    Rel
        WHERE    mult = 'single')
UNION
(SELECT Path.class, Rel.rclass
 FROM    Path, Rel
 WHERE    Path.rclass = Rel.class
        AND    Rel.mult = 'single'    )
SELECT *
FROM    Path;

```

(c) WITH RECURSIVE Path(class, rclass) AS

```

    (SELECT class, rclass
     FROM    Rel
     WHERE    mult = 'multi')
UNION
(SELECT Path.class, Rel.rclass
 FROM    Path, Rel
 WHERE    Path.rclass = Rel.class)
UNION
(SELECT Rel.class, Path.rclass
 FROM    Path, Rel
 WHERE    Rel.rclass = Path.class)
SELECT *
FROM    Path;

```

(d) This could be viewed as relation from (a) EXCEPT relation from (b).

```

WITH RECURSIVE PathAll(class, rclass) AS
    (SELECT class, rclass
     FROM    Rel    )
UNION
    (SELECT PathAll.class, Rel.rclass
     FROM    PathAll, Rel
     WHERE    PathAll.rclass = Rel.class),
RECURSIVE PathSingle(class, rclass) AS
    (SELECT class, rclass
     FROM    Rel

```

```

        WHERE mult = 'single')
UNION
(SELECT PathSingle.class, Rel.rclass
 FROM   PathSingle, Rel
 WHERE  PathSingle.rclass = Rel.class
        AND Rel.mult = 'single'      )
SELECT class, rclass
FROM   PathAll
EXCEPT
SELECT class, rclass
FROM   PathSingle
;

```

- (e) We include the edge label as part of the recursive relation and then, basically, we build the path as in (a) except we only add edges that have an opposite label.

```

WITH RECURSIVE Path(class, rclass, mult) AS
  (SELECT class, rclass, mult
   FROM   Rel
        )
UNION
  (SELECT Path.class, Rel.rclass, Rel.mult
   FROM   Path, Rel
   WHERE  Path.rclass = Rel.class
        AND Path.mult <> Rel.mult )
SELECT *
FROM   Path;

```

- (f) WITH RECURSIVE Path(class, rclass) AS
- ```

  (SELECT class, rclass
   FROM   Rel
   WHERE  mult = 'single')
UNION
  (SELECT Path.class, Rel.rclass
   FROM   Path, Rel
   WHERE  Path.rclass = Rel.class
        AND Rel.mult = 'single'      )

```

```

SELECT *
FROM   Path X
WHERE  EXISTS(SELECT 1
               FROM   Path Y
               WHERE  Y.class = X.rclass
               AND    Y.rclass = X.class )
;

```

## Section 10.3

### Exercise 10.3.1

- (a) Stars(name, address, birthdate)  
     Movies(title, year, length, stars(\*Stars))
- (b) Stars(name, address, birthdate)  
     Movies(title, year, length, stars(\*Stars))  
     Studios(name, address, movies(\*Movies))
- (c) Stars(name, address, birthdate)  
     Movies(title, year, length, studio(name, address), stars(\*Stars))

### Exercise 10.3.2

```

Customers(name, address, phone, ssNo, accts(*Accounts))
Accounts(number, type, balance, owners(*Customers))

```

### Exercise 10.3.3

```

Customers(name, address, phone, ssNo, accts(*Accounts))
Accounts(number, type, balance, owner(*Customers))

```

### Exercise 10.3.4

```

Players(name)
Teams(name, players(*Players), captain(*Players), colors)
Fans(name, fav_teams(*Teams), fav_players(*Players), fav_color)

```

### Exercise 10.3.5

People(name, mother(\*People), father(\*People), children(\*People))

## Section 10.4

### Exercise 10.4.1

```
Movies(  
  title      TitleType,  
  year       YearType,  
  length     DurationType,  
  genre      GenreType,  
  studioName BusinessNameType,  
  producerC# CertificateType  
)
```

```
MovieStar(  
  name       PersonNameType,  
  address    AddressType,  
  gender     GenderType,  
  birthdate  DateType  
)
```

```
StarsIn(  
  movieTitle TitleType,  
  movieYear  YearType,  
  starName   PersonNameType  
)
```

```
MovieExec(  
  name       PersonNameType,  
  address    AddressType,  
  cert#      CertificateType,  
  netWorth   CurrencyType  
)
```

```
Studio(  
  name       PersonNameType,  
  address    AddressType,  
  cert#      CertificateType,  
  netWorth   CurrencyType  
)
```

```

name      BusinessNameType,
address   AddressType,
presC#    CertificateType
)

```

### Exercise 10.4.2

- (a) CREATE TYPE NameType AS(
 first VARCHAR(30),
 middle VARCHAR(50),
 last VARCHAR(30),
 title VARCHAR(10)
 );
- (b) CREATE TYPE PersonType AS(
 name NameType,
 mother REF(PersonType),
 father REF(PersonType)
 );
- (c) CREATE TYPE MarriageType AS(
 date DATE,
 husband REF(PersonType),
 wife REF(PersonType)
 );

### Exercise 10.4.3

```

CREATE TYPE ProductType AS(
  maker    CHAR(5),
  model     INTEGER,
  type      CHAR(8)
);

CREATE TABLE Product OF ProductType(
  REF IS ProductId SYSTEM GENERATED
);

```

```

CREATE TABLE PC(
    model      REF(ProductType) SCOPE Product,
    speed      DECIMAL(5,2),
    ram        INTEGER,
    hd         INTEGER
    price      DECIMAL(10,2)
);

```

```

CREATE TABLE Laptop(
    model      REF(ProductType) SCOPE Product,
    speed      DECIMAL(5,2),
    ram        INTEGER,
    hd         INTEGER
    screen     DECIMAL(5,2),
    price      DECIMAL(10,2)
);

```

```

CREATE TABLE Printer(
    model      REF(ProductType) SCOPE Product,
    color      CHAR(1),
    type       VARCHAR(10),
    price      DECIMAL(10,2)
);

```

### Exercise 10.4.4

Model attribute in Products cannot be a reference to the tuple in the relation for that type of product because that would create a circular reference situation where the model is a reference to the relation itself which has a model attribute but is a reference, etc. There would not be a column that stores the actual model values.

### Exercise 10.4.5

```

CREATE TYPE ClassType AS (
    class      VARCHAR(30),
    type       CHAR(2),
    country    VARCHAR(30),
    numGuns    INTEGER,

```

```

    bore            INTEGER,
    disp            INTEGER
);

CREATE TYPE ShipType AS (
    name            VARCHAR(30),
    class           REF(ClassType),
    launched        INTEGER
);

CREATE TYPE BattleType AS (
    name            VARCHAR(30),
    date            DATE
);

CREATE TYPE OutcomeType AS (
    ship            REF(ShipType),
    battle          REF(BattleType),
    result          VARCHAR(10)
);

CREATE TABLE Classes OF ClassType (
    REF IS classID SYSTEM GENERATED
);

CREATE TABLE Ships OF ShipType(
    REF IS shipID SYSTEM GENERATED
);

CREATE TABLE Battles OF TYPE BattleType(
    REF IS battleID SYSTEM GENERATED
);

CREATE TABLE Outcomes OF TYPE OutcomeType(
    REF IS outcomeID SYSTEM GENERATED
);

```

## Section 10.5

### Exercise 10.5.1

- (a) 

```
SELECT star->name
FROM   StarsIn
WHERE  movie->title = 'Dogma';
```
- (b) 

```
SELECT DISTINCT movie->title, movie->year
FROM   StarsIn
WHERE  star->address.city() = 'Malibu';
```
- (c) 

```
SELECT movie
FROM   StarsIn
WHERE  star->name = 'Melanie Griffith';
```
- (d) 

```
SELECT  movie->title, movie->year
FROM    StarsIn
GROUP BY movie->title, movie->year
HAVING  COUNT(*) >= 5;
```

### Exercise 10.5.2

- (a) 

```
SELECT model->maker
FROM   PC
WHERE  hd > 60;
```
- (b) 

```
SELECT DISTINCT model->maker
FROM   Printers
WHERE  type = 'laser';
```
- (c) 

```
WITH MaxSpeedsPerMaker(maker, maxSpeed) AS(
    SELECT  model->maker, MAX(speed)
    FROM    Laptops
    GROUP BY model->maker
),
MakerTopModel(maker,topModel) AS(
    SELECT M.maker, L.model->model
    FROM    Laptops L, MaxSpeedsPerMaker M
    WHERE   L.model->maker = M.maker
           AND L.speed      = maxSpeed
)
```



```

SELECT model->model, topModel
FROM   Laptops L, MakerTopModel M
WHERE  L.model->maker = M.maker
;

```

### Exercise 10.5.3

- (a) 

```

SELECT x.name
FROM   Ships x
WHERE  x.class->disp > 35000;

```
- (b) 

```

SELECT DISTINCT x.battle->name
FROM   Outcomes x
WHERE  x.result = 'sunk';

```
- (c) 

```

SELECT DISTINCT x.class->class
FROM   Ships x
WHERE  x.launched > 1930;

```
- (d) 

```

SELECT DISTINCT x.battle->name
FROM   Outcomes x
WHERE  x.result = 'damaged'
      AND x.ship->class->country = 'USA';

```

### Exercise 10.5.4

```

CREATE FUNCTION StarLEG(p1 StarType,
                        p2 StarType )
RETURNS INTEGER
  IF      p1.name < p2.name THEN RETURN(-1)
  ELSEIF  p1.name > p2.name THEN RETURN( 1)
  ELSE    RETURN(AddrLEG(p1.address,p2.address))
  ENDIF
;
CREATE ORDERING FOR StarType
  ORDERING FULL BY RELATIVE WITH StarLEG;

```

### Exercise 10.5.5

```
CREATE PROCEDURE DeleteStar(IN pName VARCHAR(50))
BEGIN
    DELETE FROM StarsIn
    WHERE star->name = pName;

    DELETE FROM MovieStar x
    WHERE x.name = pName;
END;
```

## Section 10.6

### Exercise 10.6.1

- (a) Dimension attributes are: cust, date, proc, memory, hd, od.  
Dependent attributes are: quant, price.
- (b) Cust(custID, name, address, phone, creditCard)  
Proc(procID, manufacturer, name, model, speed)  
HD(hdID, manufacturer, name, model, capacity, cylinders, surfaces, speed)  
OD(odID, manufacturer, type, capacity, speed)

### Exercise 10.6.2

First we could select the number of orders that had DVD disks and the number of orders that had CD disks. This would show just the totals over all orders.

```
SELECT  D1.type, COUNT(*)
FROM    Orders F, OD D1
WHERE   F.od = D1.odID
GROUP BY D1.type
HAVING  D1.type IN('DVD', 'CD')
;
```

Then we could drill-down to see what the totals are per month, hopefully seeing that the numbers for DVDs increase and the numbers for CDs decrease.

```

SELECT    MONTH(F.date) MONTHS, D1.type, COUNT(*)
FROM      Orders F, OD D1
WHERE     F.od = D1.odID
GROUP BY  MONTHS, D1.type
HAVING    D1.type IN('DVD', 'CD')
;

```

Next we could drill-up to show the totals per year.

```

SELECT    YEAR(F.date) YEARS, D1.type, COUNT(*)
FROM      Orders F, OD D1
WHERE     F.od = D1.odID
GROUP BY  YEARS, D1.type
HAVING    D1.type IN('DVD', 'CD')
;

```

## Section 10.7

### Exercise 10.7.1

- (a) The ratio is  $\left(\frac{11}{10}\right)^{10}$ , or about 2.59.
- (b) The ratio is  $\left(\frac{3}{2}\right)^{10}$ , or about 57.66.

### Exercise 10.7.2

- (a) Assuming the column name for SUM(val) in SalesCube is val:

```

SELECT    dealer, val
FROM      SalesCube
WHERE     model IS NULL
        AND color = 'blue'
        AND date IS NULL
        AND dealer IS NOT NULL
;

```

- (b) Assuming the column name for SUM(cnt) in SalesCube is cnt:

```

SELECT    cnt
FROM      SalesCube
WHERE     model = 'Gobi'
        AND color = 'green'
        AND date IS NULL
        AND dealer = 'Smilin'' Sally'
;

```

- (c) Assuming the column names for SUM(cnt) and SUM(val) in SalesCube are cnt and val:

```

SELECT    val/cnt
FROM      SalesCube
WHERE     model = 'Gobi'
        AND color IS NULL
        AND YEAR(date) = 2007
        AND MONTH(date) = 3
        AND dealer IS NOT NULL
;

```

### Exercise 10.7.3

The rollup would not help and would make it more difficult to ensure that we do not double count the rows and only consider the rows that are in CUBE(Sales) but not in Sales.

### Exercise 10.7.4

```

CREATE MATERIALIZED VIEW OrdersCube(
    cust, date, proc, memory, hd, od, tquant, tprice)
AS(
    SELECT cust, date, proc, memory, hd, od, SUM(quant), SUM(price)
    FROM    Orders
    GROUP BY cust, date, proc, memory, hd, od)
WITH CUBE;

```

### Exercise 10.7.5

- (a) 

```
SELECT D1.speed, MONTH(F.date), SUM(F.tquant)
FROM   OrdersCube F, Proc D1
WHERE  F.proc = D1.procID
      AND F.cust IS NULL
      AND YEAR(F.date) = 2007
      AND F.memory IS NULL
      AND F.hd IS NULL,
      AND F.od IS NULL
GROUP BY D1.speed, MONTH(F.date)
;
```
- (b) 

```
SELECT D1.type, D2.type, SUM(F.tquant)
FROM   OrdersCube F, Proc D1, HD D2
WHERE  F.proc = D1.procID
      AND F.hd   = D2.hdID
      AND F.cust IS NULL
      AND F.date IS NULL
      AND F.memory IS NULL
      AND F.od IS NULL
GROUP BY D1.type, D2.type
;
```
- (c) 

```
SELECT MONTH(F.date), SUM(tprice)/SUM(F.tquant)
FROM   OrdersCube F, Proc D1
WHERE  F.proc = D1.procID
      AND D1.speed = 3.0
      AND F.cust IS NULL
      AND F.date >= '01/01/2005'
      AND F.memory IS NULL
      AND F.hd IS NULL,
      AND F.od IS NULL
GROUP BY MONTH(F.date)
;
```

### Exercise 10.7.6

Yes, other rollups could contain these tuples. Those rollups can be formed by rearranging the group by list so that columns we need to be aggregated are at the tail of the list. For instance, to include tuple

```
('Gobi', NULL, '2001-05-21', 'Friendly Fred', 152000, 7)
```

The group by list would be:

```
GROUP BY model, date, dealer, color WITH ROLLUP
```

### Exercise 10.7.7

In the worst case, the fact table could have only one row, the CUBE(F) would add an additional  $2^n$  tuples, and so the ratio would be  $2^n$ .